Image Coding Audiovisual Processing CMP-6026A

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December 01 2021

Content

Lossy and lossless image compression.

- Changing colour spaces and subsampling
- DCT and quantisation
- Run-length encoding
- Entropy coding

Image Coding

How can we compress an image without *destroying* the image?

- Data and information are not the same thing.
- Goal is to identify and remove **redundancy**.

- Image can be reconstructed **exactly**.

- Inflated image is an **approximation** of the original.
- How much loss is *acceptable*?

Image Redundancy

Inter-pixel redundancy:

- Neighbouring pixels are related to one another

Image Redundancy

Coding redundancy:

- Not all pixel intensities are equally likely

Image Redundancy

Pycho-visual redundancy:

- We are not visually *sensitive* to everything in the image

- A framework for compressing images.
- Many algorithms can be used in the framework.
- Developed by Joint Photographic Expert Group.
- JPEG exploits the three forms of redundancy outlined.



Figure 1: YC_bC_r

$$Y = 0.299R + 0.587G + 0.114B$$
$$C_b = B - Y$$
$$C_r = R - Y$$

Luminance

Y = 0.299R + 0.587G + 0.114B

Humans are *more* sensitive to luminance...

Chrominance

$$C_b = B - Y$$
$$C_r = R - Y$$

Humans are *less* sensitive to chrominance...

We can downsample the chrominance channels without affecting the image in a *perceptible* way.

- Exploits psycho-visual redundancy.



Figure 2: Chroma Subsampling

Subsampling scheme is expressed as a ratio J:a:b

- represents a conceptual window on the *chrominance* channels.

- J: horizontal sampling reference. Usually, 4.
- a: number of pixels in the top row that will have chroma information.
- b: number of *changes* of samples (Cr, Cb) between first and second row of J pixels.

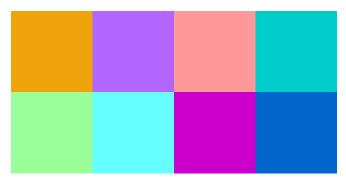


Figure 3: Chroma Subsampling



1	2	3	4
1	2	3	4

Figure 4: 4:4:4



1	2
1	2

Figure 5: 4:2:2



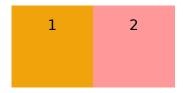


Figure 6: 4:2:0

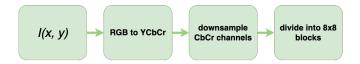


Figure 7: 8x8 Blocks

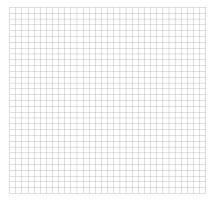


Figure 8: image matrix

block	block	block	block
1	2	3	4
block	block	block	block
5	6	7	8
block	block	block	block
9	10	11	12

Figure 9: 8x8 blocks



Figure 10: DCT

Transforms the image into the frequency domain.

DCT

8x8 Block	Intensities							
	36	90	140	141	34	135	33	32
	172	198	186	188	48	84	39	63
	201	209	204	210	151	43	114	151
	190	202	208	209	208	115	35	33
	172	183	189	203	210	171	39	34
	138	173	190	193	209	175	40	39
	114	159	181	182	200	185	49	38
	131	53	40	37	66	85	39	35

Figure 11: image values

DCT

8x8 Block	DCT Coefficients	
	1041 -297 194 6 -10 34 -48 53	1
	-20 -39 -45 121 19 -47 14 40	0
	232 101 -75 -2 50 -41 -33 4	1
	-5 23 129 -27 59 25 -48 44	4
	124 29 6 -57 6 -5 -24 0	,
	-41 22 71 10 9 19 2 -1	5
	28 53 -42 13 -21 -9 15 -2	5
	-36 39 -5 15 -5 -3 9 -6	5

Figure 12: coefficients

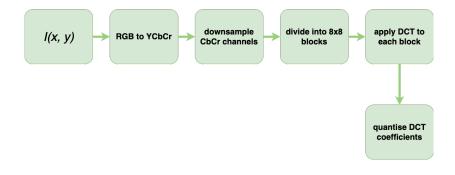


Figure 13: DCT Quantisation

Reduce the number of bits needed to store a value by reducing precision.

- Decrease precision as we move away from the top left corner.
- High frequency details usually contribute less to the image.

Quantisation is performed as follows:

$$DCT_q(i,j) = round\left(\frac{DCT(i,j)}{Q(i,j)}\right)$$

where Q is the quantisation matrix.

	Quantisation Matrix Q												
16	11	10	16	24	40	51	61						
12	12	14	19	26	58	60	55						
14	13	16	24	40	57	69	56						
14	17	22	29	51	87	80	62						
18	22	37	56	68	109	103	77						
24	35	55	64	81	104	113	92						
49	64	78	87	103	121	120	101						
72	92	95	98	112	100	103	99						

Figure 14: quantisation matrix

DCT Coefficients Q									Qua	antis	ed [ОСТ												
1041	-297	194	6	-10	34	-48	51	16	11	10	16	24	40	51	61		65	-27	19	0	0	1	-1	1
-20	-39	-45	121	19	-47	14	40	12	12	14	19	26	58	60	55		-2	-3	-3	6	1	-1	0	1
232	101	-75	-2	50	-41	-33	41	14	13	16	24	40	57	69	56		17	8	-5	0	1	-1	0	1
-5	23	129	-27	59	25	-48	44	14	17	22	29	51	87	80	62		0	1	6	-1	1	0	-1	1
124	29	6	-57	6	-5	-24	0	18	22	37	56	68	109	103	77		7	1	0	-1	0	0	0	0
-41	22	71	10	9	19	2	-15	24	35	55	64	81	104	113	92		-2	1	1	0	0	0	0	0
28	53	-42	13	-21	-9	15	-25	49	64	78	87	103	121	120	101		1	1	-1	0	0	0	0	0
-36	39	-5	15	-5	-3	9	-6	72	92	95	98	112	100	103	99		-1	0	0	0	0	0	0	0

Figure 15: quantisation

_	Quantised DCT													
65	-27	19	0	0	1	-1	1							
-2	-3	-3	-3 6 1 -		-1	0	1							
17	8	-5	0	1	-1	0	1							
0	1	6	6 -1 1 0		0	-1	1							
7	1	0	-1	0	0	0	0							
-2	1	1	0	0	0	0	0							
1	1	-1	0	0		0	0							
-1	0	0	0	0	0	0	0							

Figure 16: DCT_q

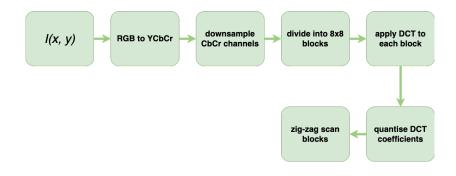


Figure 17: DCT Quantisation

ZigZag Scan

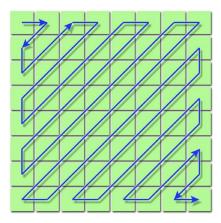


Figure 18: ZigZag Scan

ZigZag Scan

	Quantised DCT													
65	-27	19	0	0	1	-1	1							
-2	-3	-3	-3 6		-1	0	1							
17	8	-5	0	1	-1	0	1							
0	1	6	-1	1	0	-1	1							
7	1	0	-1	0	0	0	0							
-2	1	1	0	0	0	0	0							
1	1	-1	0	0	0	0	0							
-1	0	0	0	0	0	0	0							

Figure 19: quantised block

ZigZag Scan

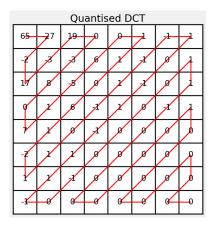


Figure 20: ZigZag Scan

$$65, -27, -2, 17, -3,$$

 $19, 0, -3, 8, 0, \dots$

ZigZag Scan

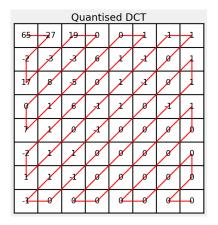


Figure 21: ZigZag Scan

Reads from low frequency coefficients to high frequency coefficients...

ZigZag Scan

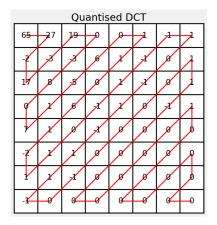


Figure 22: ZigZag Scan

More likely to encode all non-zeros and all zeros together...

beneficial for the next step...

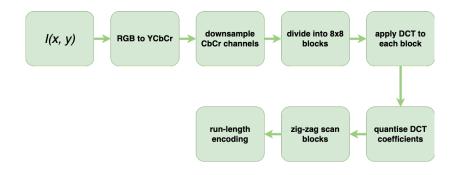


Figure 23: run-length encoding

Extracts series of value and length of runs from sequence of values. Exploits **inter-pixel** redundancy.

65 -27 -2 17 -3 -3 1 1 1 -2 1 1 0 -1 1 **0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0**

Exploits inter-pixel redundancy

 the relationship between neighbouring "pixels" in the zigzag scan of the DCT coefficient matrix

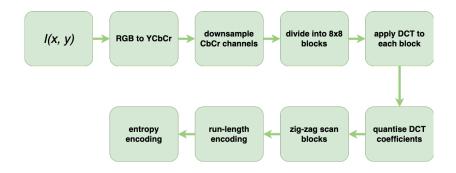


Figure 24: entropy encoding

Entropy Coding

Information and data are not the same thing.

 Claude Shannon, (1948). A Mathematical Theory of Communication. Entropy Coding exploits coding redundancy

- not every value is equally likely.

Entropy Coding encodes a sequence with *variable* length code so that:

- $-\,$ More probable values have fewer bits, and
- less probable values have more bits.

The new alphabet requires fewer bits per pixel.

- How many bits do we need?

Recall: the *probability* of an event is:

$$p_i = \frac{N_i}{N}$$

The *information* in **bits** is:

$$I_i = -\log_2 p_i$$

The entropy, the smallest possible mean symbol length, is:

$$H = -\sum_i p_i \log_2 p_i$$

We can use these properties to develop a better coding for an image.

- The stream must be decoded *unambiguously*.
- One code cannot be the **prefix** of another.

Step 1:

- Arrange values in order of decreasing probability.
- Each forms a *leaf* in the **Huffman tree**.

Step 2:

- Merge the two leaves with the smallest probability,
 - add the probabilities,
 - insert the node into the sorted list.
- Assign a 1/0 to each branch being merged.

Step 3:

- Repeat until only the root node remains.
- $-\,$ Read codewords from the root to the leaves.

0	0	1	2	5	5	7	4	5	5
5	1	1	2	1	4	1	4	3	1
5	2	1	2	1	2	2	5	3	1
3	7	2	6	5	3	5	5	1	1
2	7	5	4	5	5	5	3	1	1
7	4	5	5	5	5	5	3	3	5
1	5	5	5	5	1	1	2	2	5
6	5	7	4	2	1	4	1	2	5
1	1	7	2	1	2	4	1	3	5
1	2	0	0	7	4	7	7	4	5

What is the Huffman code for this image? And, what is the current bit rate? Count the frequencies of each symbol.

Frequency	Symbol
4	0
23	1
15	2
8	3
10	4
29	5
2	6
9	7

What is the **entropy** of this image?

p(s)	$-\log p(s)$	×
0.29	1.786	0.518
0.23	2.120	0.488
0.15	2.737	0.411
0.10	3.322	0.332
0.09	3.474	0.313
0.08	3.644	0.292
0.04	4.644	0.186
0.02	5.644	0.113
	+	2.651

Sort by the most frequent symbol.

Frequency	Symbol
29	5
23	1
15	2
10	4
9	7
8	3
4	0
2	6

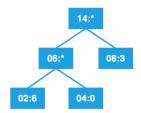
Merge the two leaves with the lowest frequency...



Insert the node into the sorted list.

Frequency	Symbol
29	5
23	1
15	2
10	4
9	7
8	3 *
6	*

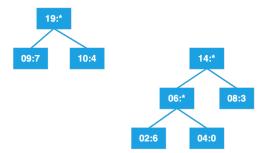
Repeat with the next two lowest frequencies.



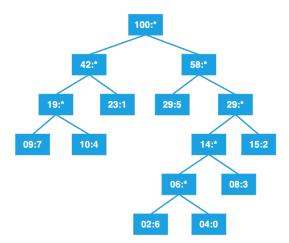
Frequency	Symbol
29	5
23	1
15	2 *
14	*
10	4
9	7

Insert the node into the sorted list.

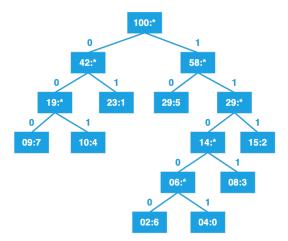
Repeat with the next two lowest frequencies.



Continue until the tree is complete.



Label left branches with **0**, right branches with **1**.



Read from the **root** to compute the new codes.

Code	Symbol
11001	0
01	1
111	2
1101	3
001	4
10	5
11000	6
000	7

Value p(x) code length 5 0.29 2 1 0.23 2 2 0.15 3 4 0.10 3 7 0.09 3 3 0.08 4	× 0.58
1 0.23 2 2 0.15 3 4 0.10 3 7 0.09 3	0.58
2 0.15 3 4 0.10 3 7 0.09 3	
4 0.10 3 7 0.09 3	0.46
7 0.09 3	0.45
	0.30
2 0.09 1	0.27
5 0.00 4	0.32
0 0.04 5	0.20
6 0.02 5	0.10
+	2.68

We can calculate the bit rate we achieved.

- Not optimal.
- optimal bit rate is 2.65
- our bit rate is 2.68
- The compression ratio is 2.68/3.0 = 0.8933.

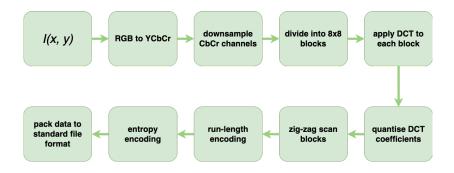


Figure 25: data packing

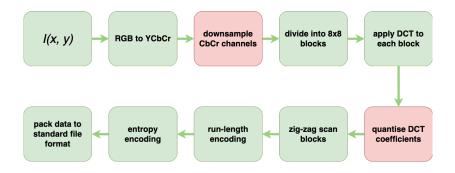


Figure 26: lossy components



Figure 27: 50% quality



Figure 28: 5% quality

Summary

Three types of redundancy are exploited in image compression.

- psycho-visual redundancy
- inter-pixel redundancy
- coding redundancy
- JPEG uses them all.